

Multidimensional Dataflow

Ptolemy ILP Conference

March 14, 1997

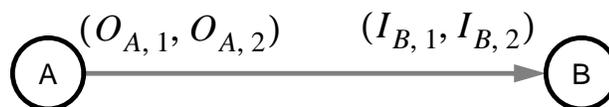
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Multidimensional Dataflow Extension



Balance equations:

$$r_{A,1} O_{A,1} = r_{B,1} I_{B,1}$$

$$r_{A,2} O_{A,2} = r_{B,2} I_{B,2}$$

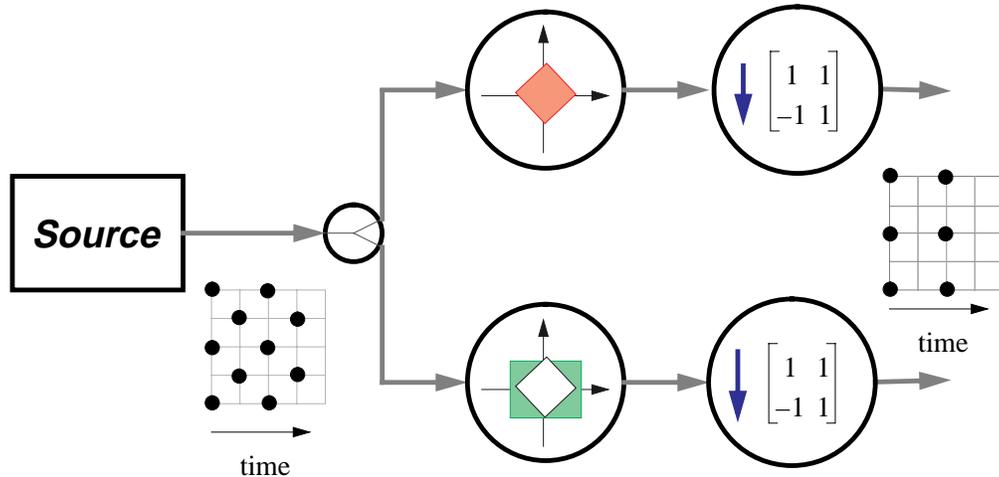
Solve for the smallest integers $r_{X,i}$, which then give the number of repetitions of actor X in dimension i .

Higher dimensionality follows similarly.

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Generalization to Arbitrary Lattices

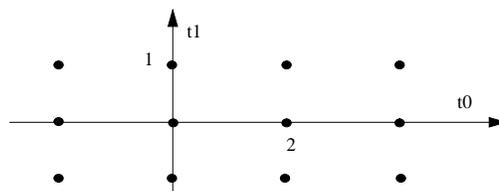
- MDSDF handles only rectangularly sampled signals.
- GMDSDF handles signals on arbitrary lattices, *without sacrificing compile-time schedulability.*



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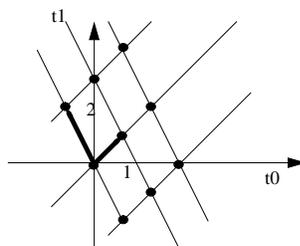
Non-rectangular Sampling

Rectangular sampling



$$V = \begin{bmatrix} 2 & 0 \\ 0 & 1 \end{bmatrix}$$

Non-rectangular sampling



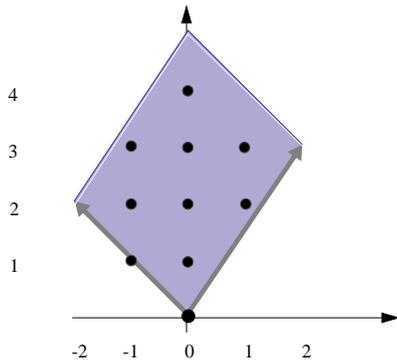
$$V = \begin{bmatrix} 1 & -1 \\ 1 & 2 \end{bmatrix}$$

Definition: The set of all sample points given by $\hat{t} = V\hat{n}$, $\hat{n} \in \mathfrak{N}$ is called the *lattice* generated by V . It is denoted $LAT(V)$.

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The Fundamental Parallelepiped

The *fundamental parallelepiped*, denoted by $FPD(V)$, is the set of points given by Vx where $x = [x_1, x_2]^T$ with $0 \leq x_1, x_2 < 1$.



$$L = \begin{bmatrix} 2 & -2 \\ 3 & 2 \end{bmatrix}$$

Definition: The set of integer points in $FPD(V)$ is denoted as $N(V)$.

Lemma: $J(V) = |N(V)| = |\det(V)|$ for an integer matrix V .

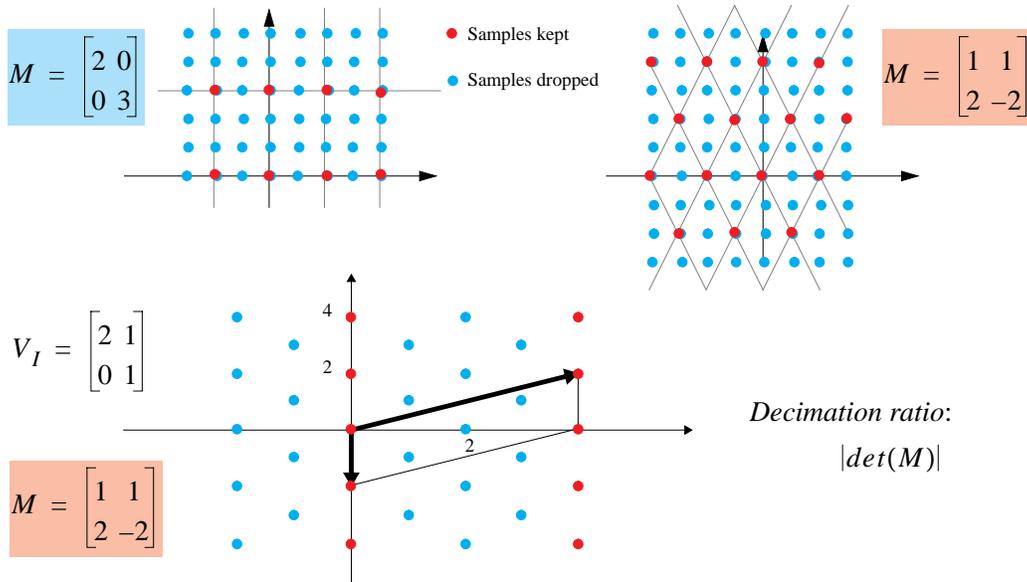
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Multidimensional Decimators

M-D decimation is given by the relationship:

$$y(\hat{n}) = x(\hat{n}), \hat{n} \in LAT(V_I M)$$

where x is defined on the points $V_I k$, V_I being the sampling matrix.



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Multidimensional Expanders

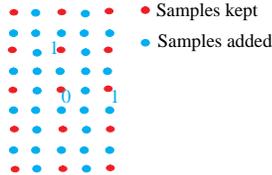
M-D expander:

$$y(n) = \begin{cases} x(n) & n \in \text{LAT}(V_I) \\ 0 & \text{otherwise} \end{cases} \forall n \in \text{LAT}(V_I L^{-1})$$

where x is defined at the points $V_I k$, V_I being the sampling matrix.

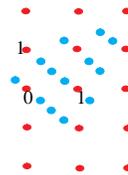
Rectangular expansion

$$L = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix}$$



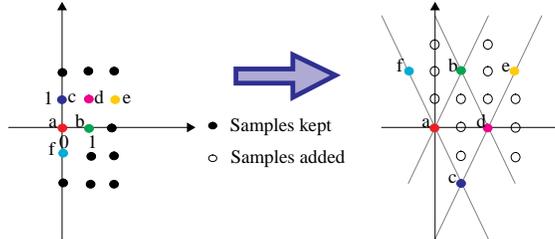
Non-rectangular expansion

$$L = \begin{bmatrix} 1 & 1 \\ 2 & -2 \end{bmatrix}$$



Renumbered samples from the expanders output

$$L = \begin{bmatrix} 1 & 1 \\ 2 & -2 \end{bmatrix}$$



Expansion ratio:

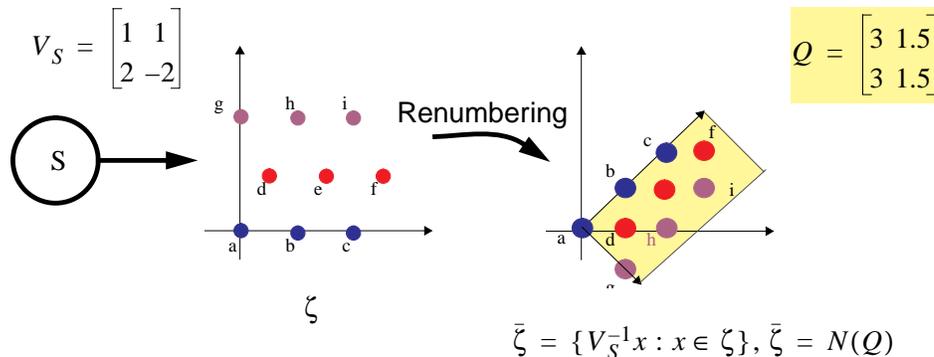
$$|\det(L)|$$

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Generalized MDSDF (GMDSDF): Sources

Definition: The **containability condition:** let X be a set of integer points in \mathfrak{R}^m . We say that X satisfies the *containability condition* if there exists an $m \times m$ matrix W such that $N(W) = X$.

Definition: We will assume that any source actor in the system produces data in the following manner. A source S will produce a set of samples ζ on each firing such that each sample in ζ will lie on the lattice $\text{LAT}(V_S)$. We assume that the renumbered set $\bar{\zeta}$ satisfies the containability condition.

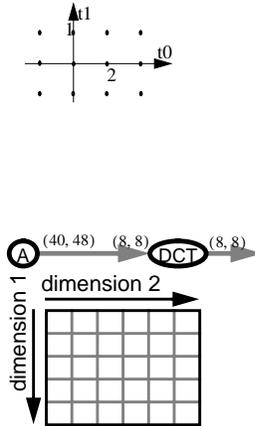


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Concise Problem Statement

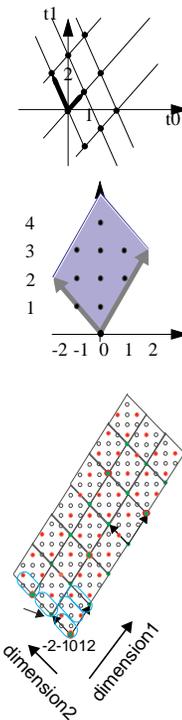
MDSDF

- Rectangular lattice
- Regions of data produced = rectangular arrays
- Rectangular arrays specified concisely by tuples of produced/ consumed.
- Coordinate axes for dataflow along arcs orthogonal to each other (x and y axes).



GMDSDF

- Arbitrary lattice
- Regions of data produced = parallelograms
- Parallelograms specified concisely as the set of integer points inside a support matrix.
- Coordinate axes for dataflow along arcs not necessarily orthogonal.



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Support Matrices

Want to describe regions where the data is contained.

- In MDSDF, these are ordinary arrays
- In the extension, these are *support matrices*.



Theorem:

For the decimator,

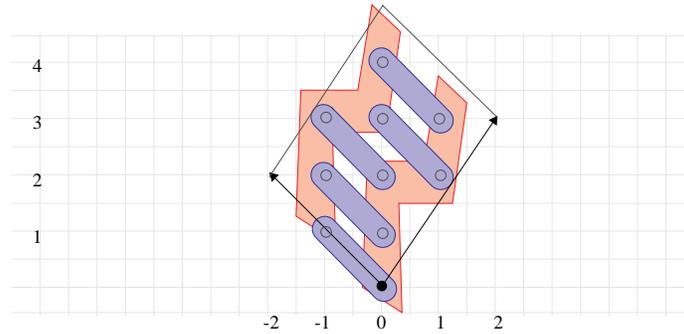
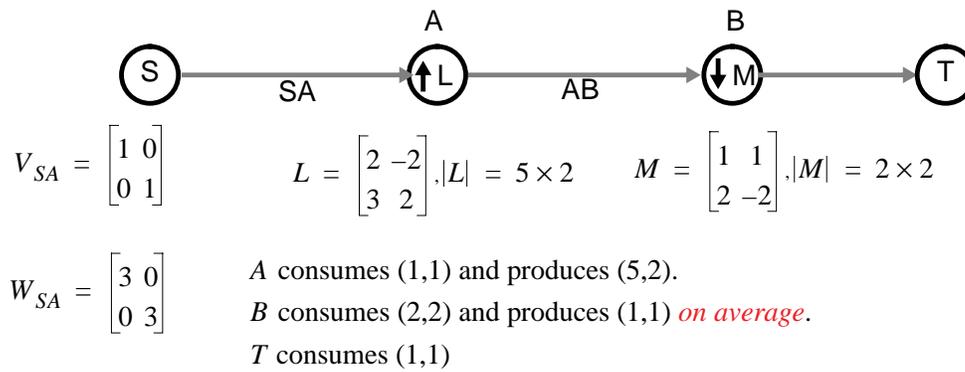
$$V_f = V_e M \text{ and } W_f = M^{-1} W_e.$$

For the expander,

$$V_f = V_e L^{-1}, \text{ and } W_f = L W_e.$$

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Semantics of GMDSDF



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GMDSDF — Balance Equations

- We don't know yet exactly how many samples on each firing the decimator will produce.
- Idea: Assume that it produces (1,1) and solve balance equations:

$$3r_{S,1} = 1r_{A,1} \quad 5r_{A,1} = 2r_{B,1} \quad r_{B,1} = r_{T,1}$$

$$3r_{S,2} = 1r_{A,2} \quad 2r_{A,2} = 2r_{B,2} \quad r_{B,2} = r_{T,2}$$

- Solution:

$$r_{S,1} = 2, r_{S,2} = 1$$

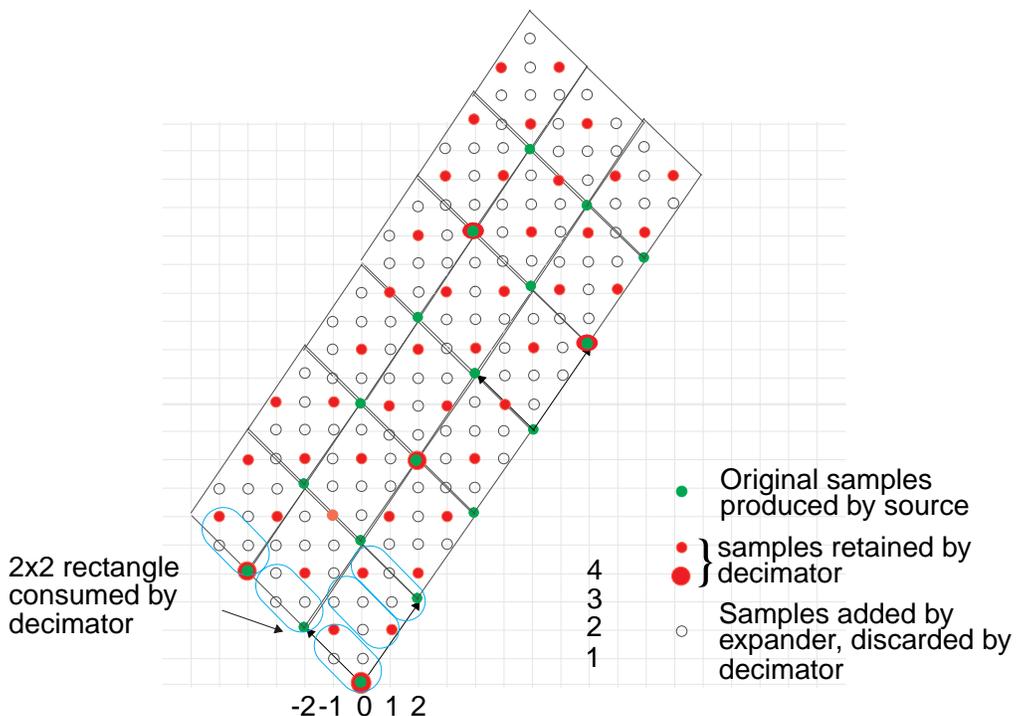
$$r_{A,1} = 6, r_{A,2} = 3$$

$$r_{B,1} = 15, r_{B,2} = 3$$

$$r_{T,1} = 15, r_{T,2} = 3$$

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Dataspace on arc AB



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Balance equations cont'd

Question: Have we really “balanced”?

No: by counting the number of samples that have been kept in the previous slide.

More systematically:

$$W_{SA} = \begin{bmatrix} 3 & 0 \\ 0 & 3 \end{bmatrix} \begin{bmatrix} r_{S,1} & 0 \\ 0 & r_{S,2} \end{bmatrix} = \begin{bmatrix} 3r_{S,1} & 0 \\ 0 & 3r_{S,2} \end{bmatrix}$$

$$W_{AB} = LW_{SA} = \begin{bmatrix} 2 & -2 \\ 3 & 2 \end{bmatrix} \begin{bmatrix} 3r_{S,1} & 0 \\ 0 & 3r_{S,2} \end{bmatrix} = \begin{bmatrix} 6r_{S,1} & -6r_{S,2} \\ 9r_{S,1} & 6r_{S,2} \end{bmatrix}$$

$$W_{BT} = M^{-1}W_{AB} = \frac{1}{4} \begin{bmatrix} 2 & 1 \\ 2 & -1 \end{bmatrix} \begin{bmatrix} 6r_{S,1} & -6r_{S,2} \\ 9r_{S,1} & 6r_{S,2} \end{bmatrix} = \frac{1}{4} \begin{bmatrix} 21r_{S,1} & -6r_{S,2} \\ 3r_{S,1} & -18r_{S,2} \end{bmatrix}$$

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Balance equations cont'd

Want to know if

$$|N(W_{BT})| = \frac{|N(W_{AB})|}{|M|}$$

We have

$$|N(W_{AB})| = |\det(W_{AB})| = 90r_{S,1}r_{S,2}$$

The right hand side becomes

$$\frac{90r_{S,1}r_{S,2}}{4} = \frac{45r_{S,1}r_{S,2}}{2}$$

Therefore, we need

$$r_{S,1}r_{S,2} = 2k \quad k = 0, 1, 2, \dots$$

The balance equations gave us $r_{S,1} = 2, r_{S,2} = 1$.

With these values, we get

$$W_{BT} = \begin{bmatrix} 21/2 & -3/2 \\ 3/2 & -9/2 \end{bmatrix}.$$

This matrix has 47 points inside its FPD (determined by drawing it out).

==> Balance equation solution is not quite right.

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Augmented Balance Equations

To get the correct balance, take into account the constraint given by

$$|N(W_{BT})| = \frac{|N(W_{AB})|}{|M|}$$

Sufficiency: force W_{BT} to be an integer matrix.

$$\implies r_{S,1} = 4k, k = 1, 2, \dots$$

$$\implies r_{S,2} = 2k, k = 1, 2, \dots$$

Therefore,

$$r_{S,1} = 4, r_{S,2} = 2.$$

- So decimator produces (1,1) on average but has cyclostatic behavior.

Production sequence: 2,1,1,2,1,0,1,1,0,1,2,1,1,2,1,...

Theorem:

Always possible to solve these *augmented* balance equations.

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Effect of Different Factorizations

Suppose we let $|\det(M)| = 1 \times 4$ instead. Balance equations give:

$$r_{S,1} = 1, r_{S,2} = 2$$

$$r_{A,1} = 3, r_{A,2} = 6$$

$$r_{B,1} = 15, r_{B,2} = 3$$

$$r_{T,1} = 15, r_{T,2} = 3$$

Also,

$$W_{BT} = \begin{bmatrix} 21/4 & -3 \\ 3/4 & -9 \end{bmatrix}$$

It turns out that

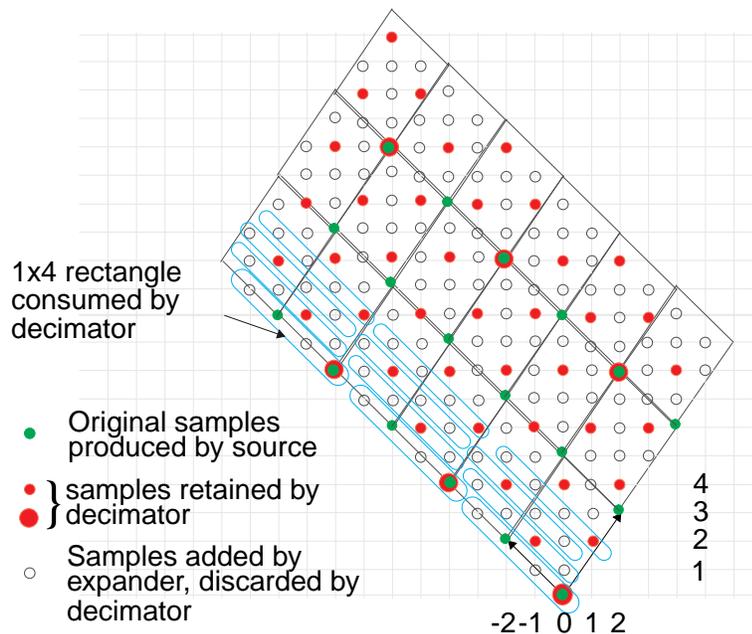
$$|N(W_{BT})| = 45$$

as required.

==> Lower number of overall repetitions with this factoring choice.

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Dataspace on Arc AB



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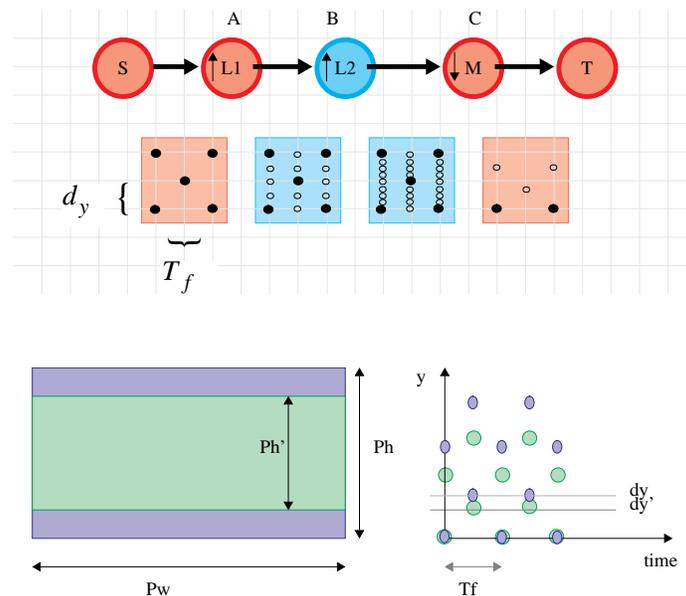
Summary of Extended Model

- Each arc has associated with it a lattice-generating matrix, and a support matrix.
- The source actor for an arc establishes the ordering of the data on that arc.
- Expander: consumes (1,1) and produces $FPD(L)$, ordered as an (L_1, L_2) rectangle where $L_1 L_2 = |det(L)|$.
- Decimator: consumes an (M_1, M_2) rectangle, where $M_1 M_2 = |det(M)|$ and produces (1,1) on average.
- Write down balance equations.
- Additional equations for support matrices on decimator outputs.
- The above two sets are simultaneously solved to determine the smallest non-zero number of times each node is to be invoked in a periodic schedule.
- Actors are then scheduled as in SDF or MDSDF.

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Aspect Ratio Conversion

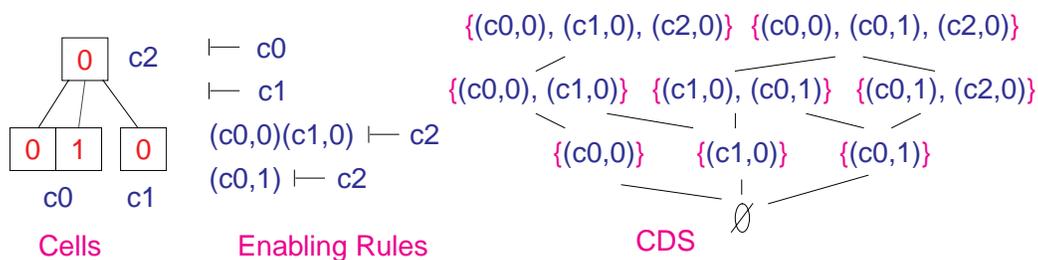
Format conversion of 2:1 interlaced video from 4/3 aspect ratio to 16/9 aspect ratio.



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Concrete Data Structures

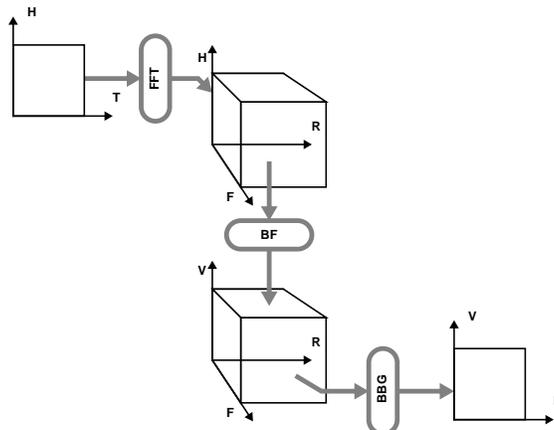
- “Cells” can have specific “Values”
- Enabling relationship says when a cell can be filled.
- “Cell” dependency partial order can be arbitrary
- Formalizes most forms of “real-world” data structures: lists, trees, arrays etc.
- Kahn-Plotkin sequential functions on CDS provide an elegant model of computation with many formal properties, like full abstraction.
- CDS approach has been mostly semantic; need to sort out operational issues (like scheduling).



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Array-OL

- Array-oriented language developed at Thomson
- Graphical syntax for specifying “array access patterns”
 - In many multidimensional programs, manipulating data aligned in various dimensions is a challenge. For example: **Transpose**.
 - Patterns specified by “paving” and “tiling” relationships.
- Combine with MDSDF...



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